

# 리드솔로몬 복호기에서 2개의 오류시, 오류위치를 찾는 최적화 방법

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## Optimizing the Circuit for Finding 2 Error Positions of 2 Error Correcting Reed Solomon Decoder

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요 약

본 논문에선 리드 솔로몬 복호기의 2개의 8비트 심볼 오류정정회로의 에러위치추적기에대한 새로운 설계법을 제시한다. 본 설계법을 통해 기존보다 빠르고 훨씬 회로량이 줄어든 적화된 2개의 8 비트심볼 오류위치 추적기를 설계할수 있었다. 이 리드솔로몬 복호기는 거의 모든 디지털 통신및 가전기기의 데이터 보존장치로 사용되질수 있다. 특히 8비트 동작을 4비트동작으로 분화시켜 복호기의 최적화를 이뤘다.

**Key Words:** Reed-Solomon(RS), Decoder,  $GF(2^4)$ , Communication, Digital, Error location, Symbol,  $GF(2^8)$

### ABSTRACT

In this paper, we show new method to find error locations of 2 eight bit symbol errors for 2 error correcting Reed-Solomon decoder. New design is much faster and has much simpler logic circuit than the former design method. This optimization was possible by partitioning the 8 bit operations into 4 bit arithmetic and logic operations. This Reed Solomon decoder can be used for data protection of almost all digital communication and consumer electronic devices.

### I. Introduction

Nowadays almost all Digital electronic devices use Reed Solomon(RS) codec to protect their data. RS encoder is used to channelcode the source data, and is used in transmitter(for communication devices) and recorder(for Digital Audio/Video devices). RS decoder is used to correct the detected error data in receiver( for communication devices) and playback device(for Digital Audio/Video devices)<sup>[2]</sup>.

In this paper, we propose new method to optimize the 2 eight bit symbol error correcting circuit of RS decoder. Almost all RS decoder of current digital

devices use 2 ~ 4 eight bit symbol error correcting circuit in the decoder. RS decoder consists of error location finding circuit and error value evaluating circuit. To find error location in the RS codeword, we should solve nonlinear error locator polynomial. It is not easy to solve the nonlinear error locating equation. On the other hand, it is straightforward and easy to find the error values once their locations are found, since the error values can be found by solving the linear systems of equations<sup>[1,6]</sup>. So the method presented here is showing how to get the solution of non linear error locator polynomial for 2 eight bit symbol error case. To optimize the error

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location finding circuit, we transform GF(2<sup>8</sup>) galois elements to GF(2<sup>4</sup>) elements<sup>[5]</sup>, since GF(2<sup>4</sup>) operation is much simpler than GF(2<sup>8</sup>) element manipulation. Table 1 shows the GF(2<sup>4</sup>) elements table, whose primitive polynomial is X<sup>4</sup>+X<sup>3</sup>+1=0<sup>[8]</sup>. In section 2 we present the old method to get 2 error locations in GF(2<sup>8</sup>) field for 2 symbol error correcting RS decoder<sup>[7]</sup>. In section 3, we present the new method and in section 4, we compared the 2 methods so finding new method is much more optimized than the old method. Finally in section 5 we make our onclusions and also show our future works to make better improvements.

Table 1 GF(2<sup>4</sup>) elements table  
표 1. GF(2<sup>4</sup>) 원소 표

GF(2 <sup>4</sup> ) elements	code(z <sub>0</sub> z <sub>1</sub> z <sub>2</sub> z <sub>3</sub> )
0	0000
1(=α <sup>15</sup> )	1000
α	0100
α <sup>2</sup>	0010
α <sup>3</sup>	0001
α <sup>4</sup>	1001
α <sup>5</sup>	1101
α <sup>6</sup>	1111
α <sup>7</sup>	1110
α <sup>8</sup>	0111
α <sup>9</sup>	1010
α <sup>10</sup>	0101
α <sup>11</sup>	1011
α <sup>12</sup>	1100
α <sup>13</sup>	0110
α <sup>14</sup>	0011

From table 1, for example, make α<sup>11</sup> • α<sup>13</sup>=α<sup>24</sup>=α<sup>9</sup>(∵α<sup>15</sup>)=1010. Also α<sup>4</sup>+α<sup>5</sup>=0100=α.

## II. Classical Method to find 2 Error positions of 2 Error correcting Reed Solomon decoder

The error locator polynomial for 2 error case is

$$\sigma(x) = x^2 + \sigma_1 x + \sigma_2 = 0. \quad (1)$$

From Newton's identities

$$\begin{bmatrix} \sigma_1 \\ \sigma_2 \end{bmatrix} = \begin{bmatrix} S_1 & S_2 \\ S_2 & S_3 \end{bmatrix} \begin{bmatrix} S_3 \\ S_4 \end{bmatrix} \quad (2)$$

So from (2),

$$\begin{aligned} \sigma_1 &= \frac{S_2 S_3 + S_1 S_4}{S_2^2 + S_1 S_3} \text{ and} \\ \sigma_2 &= \frac{S_3^2 + S_4 S_2}{S_1 S_3 + S_2^2} \end{aligned} \quad (3)$$

Now if two error positions are α<sup>i</sup>, α<sup>j</sup> then

$$\begin{aligned} \sigma_1^2 / \sigma_2 &= \frac{(S_1 S_4 + S_2 S_3)^2}{(S_2^2 + S_1 S_3)(S_3^2 + S_4 S_2)} \\ &= (\alpha^{2i} + \alpha^{2j}) / \alpha^{i+j} \\ &= \alpha^{i-j} + \alpha^{j-i} \end{aligned} \quad (4)$$

Using equation (4), we construct the ROM which maps α<sup>i</sup>+α<sup>i</sup> → α<sup>i</sup> ∈ GF(2<sup>8</sup>) as in the table 2.

After finding σ<sub>1</sub><sup>2</sup>/σ<sub>2</sub> value by equation (4), we get α<sup>i-j</sup> from the ROM data in table 2. Then 2 error positions are as in equation (5),(6).

$$\alpha^i = (\sigma_2 \cdot \alpha^{i-j})^{\frac{1}{2}} \quad (5)$$

Table 2. α<sup>k</sup>+α<sup>k</sup> to α<sup>k</sup> Mapping ROM  
표 2. α<sup>k</sup>+α<sup>k</sup>를 α<sup>k</sup>로 변환하는 리드솔로몬 메모리

Address(α <sup>k</sup> +α <sup>k</sup> )	Data (α <sup>k</sup> )
α <sup>49</sup>	α <sup>1</sup>
α <sup>98</sup>	α <sup>2</sup>
α <sup>188</sup>	α <sup>3</sup>
α <sup>196</sup>	α <sup>4</sup>
•	•
•	•
•	•
•	•
•	•
α <sup>k</sup> +α <sup>k</sup>	α <sup>k</sup>
•	•
•	•
•	•
•	•
0	α <sup>128</sup>

$$\alpha^j = (\sigma_1 + \alpha^i) \quad (6)$$

Also the 2 error values are as in equation (7),(8).

$$E_j = (\alpha^i S_0 + S_1) / \sigma_1 \quad (7)$$

$$E_i = S_0 + E_j \quad (8)$$

〈Example〉

In GF(2<sup>8</sup>), when the code polynomial C(x)=0 is transmitted to the receiver, the received polynomial is

$$r(x) = \alpha^3 x^2 + \alpha x^4 \quad (9)$$

So two error values and locations are  $\alpha^3$  at  $x^2$  and  $\alpha$  at  $x^4$ . Verify these values.

〈sol〉

First calculate syndromes  $S_i (i = 1, 2, 3, 4)$ .

$$\begin{aligned} S_1 &= r(\alpha) = 0 \\ S_2 &= r(\alpha^2) = \alpha^7 + \alpha^9 = \alpha^{57} \\ S_3 &= r(\alpha^3) = \alpha^9 + \alpha^{13} = \alpha^{109} \\ S_4 &= r(\alpha^4) = \alpha^{11} + \alpha^{17} = \alpha^{202} \end{aligned} \quad (10)$$

$$\begin{aligned} \sigma_1^2 / \sigma_2 &= \frac{(S_1 S_4 + S_2 S_3)^2}{(S_2^2 + S_1 S_3)(S_3^2 + S_4 S_2)} \\ &= (\alpha^{166})^2 / \alpha^{114} (\alpha^{218} + \alpha^{259}) \\ &= \alpha^{77} / (\alpha^{77} + \alpha^{118}) \\ &= \alpha^{77} / \alpha^{234} = \alpha^{98} \end{aligned} \quad (11)$$

From table 2, we see  $\alpha^{ij} = \alpha^2$ . So by (5), we find

$$\begin{aligned} \alpha^i &= (\sigma_2 \cdot \alpha^{i-j})^{\frac{1}{2}} = \left( \frac{S_3^2 + S_4 S_2}{S_1 S_3 + S_2^2} \right)^{\frac{1}{2}} \cdot (\alpha^2)^{\frac{1}{2}} \\ &= (\alpha^{106} + \alpha^{147})^{\frac{1}{2}} \\ &= \alpha^4 \text{ and } \alpha^j = \sigma_1 + \alpha^i \\ &= \frac{S_3}{S_2} + \alpha^4 = \alpha^{52} + \alpha^4 = (\alpha^2 + \alpha^4) + \alpha^4 = \alpha^2 \end{aligned}$$

2 Error values are from equation (7) and (8),

$$\begin{aligned} E_2 &= (\alpha^3 + \alpha) \alpha^2 / \alpha^{52} = \alpha^{53} / \alpha^{52} = \alpha \text{ and} \\ E_1 &= \alpha^{51} + \alpha = \alpha^3. \end{aligned}$$

All these are correct.

### III. New Optimized Method to Find 2 Error Positions.

If there are 2 Errors in the code  $\in GF(2^8)$ , error locator polynomial is as equation (1). Now define  $X_1 = x / \sigma_1$ , and  $K = \sigma_2 / (\sigma_1^2)$  then error locator polynomial becomes as equation (12).

$$X_1^2 + X_1 + K = 0 \quad (12)$$

Now convert (12) to GF(2<sup>4</sup>) equation as follows.

$$\begin{aligned} X_1 &= Y_1 + \beta Y_2, \quad K = K_1 + \beta K_2 \\ \text{where } Y_1 \text{ \& } Y_2 \text{ and } K_1 \text{ \& } K_2 &\in GF(2^4). \end{aligned} \quad (13)$$

Substitute (13) to (12), then

$$\begin{aligned} Y_1^2 + \beta Y_2^2 + Y_1 + K_1 + \beta(Y_2^2 + Y_2 + K_2) &= 0. \\ (\because \beta^2 + \beta = \beta[2]) \end{aligned} \quad (14)$$

From (18), we see

$$\begin{aligned} Y_1^2 + \beta Y_2^2 + Y_1 + K_1 &= 0 \text{ and} \\ Y_2^2 + Y_2 + K_2 &= 0. \end{aligned} \quad (15)$$

This  $Y_2$  can be found using table 1.

Table 3 is made easily because there are only 16 cases of  $Y_2$  for each  $K_2$ . the 2 error positions has also following relation. If  $Y_2$  is a solution to (15),  $Y_2 + 1$  is also a solution to (15) and easily verified by substitution to equation (15) itself. Once  $Y_2$  is found,  $Y_1$  can also be found from equation (15). From (15),

$$\begin{aligned} Y_1^2 + Y_1 &= K'_1 \text{ and} \\ K'_1 &= K_1 + \beta Y_2^2 \end{aligned} \quad (16)$$

Since  $Y_2$  and  $K_1$  is already known,  $K'_1$  is also known value so equation (16) is easily solved using the ROM shown in table 3. Fig. 1 shows  $K_1$  to  $K'_1$  converter circuit and  $Y_1$  &  $Y_2$  to the final error

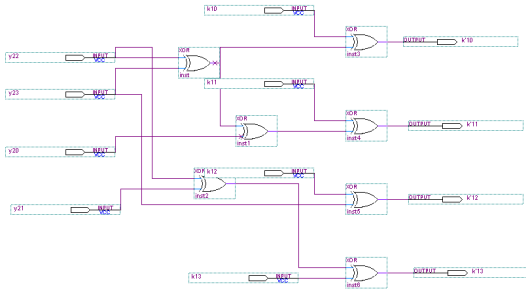


Fig. 1.  $K_1$  to  $K'_1$  Converter, where  $K_1=(K_{10}, K_{11}, K_{12}, K_{13})$ ,  $K'_1=(K'_{10}, K'_{11}, K'_{12}, K'_{13})$  and  $Y_2=(Y_{20}, Y_{21}, Y_{22}, Y_{23})$ .

그림 1.  $K_1$ 에서  $K'_1$ 으로 바꾸는 회로

Table 3.  $GF(2^4)$  2 error position look up ROM

표 3.  $GF(2^4)$  2개의 오류위치 추적 룬

$K_2=y_i^2+y_i$	$y_i$
0	0
1	$\alpha^{10}$
$\alpha$	$\alpha^5$
$\alpha^2$	No Answer
$\alpha^3$	No Answer
$\alpha^4$	No Answer
$\alpha^5$	$\alpha^7, \alpha^{13}$
$\alpha^6$	No Answer
$\alpha^7$	$\alpha^3, \alpha^4$
$\alpha^8$	No Answer
$\alpha^9$	No Answer
$\alpha^{10}$	$\alpha^{11}, \alpha^{14}$
$\alpha^{11}$	$\alpha^2, \alpha^9$
$\alpha^{12}$	No Answer
$\alpha^{13}$	$\alpha, \alpha^{12}$
$\alpha^{14}$	$\alpha^6, \alpha^8$

location finding circuit. Notice one of 2 solutions of  $Y_2$  equation (15) is invalid and can be determined by the  $K'_1$  and the ROM in table 1. If  $K'_1 \neq$  one of  $(\alpha^5, \alpha^7, \alpha^{10}, \alpha^{11}, \alpha^{13}, \alpha^{14})$ , corresponding  $Y_2$  is invalid solution<sup>[2]</sup>. If we want to find  $K'_1$  to  $Y_1$  converter, using primitive polynomial in  $GF(2^4)$ ,  $\gamma^4+\gamma^3+1=0$ , we get logic equation 17.

$$Y_1=(y_{10}, y_{11}, y_{12}, y_{13}) \\ = [K_0, K_1, K_2, K_3] = \begin{bmatrix} 0101 \\ 0000 \\ 0100 \\ 0011 \end{bmatrix} \quad (17)$$

where  $K'_1= [K_0, K_1, K_2, K_3]$ . This logic is also much simpler than the  $GF(2^8)$  case<sup>[7]</sup>. Fig. 2 shows the

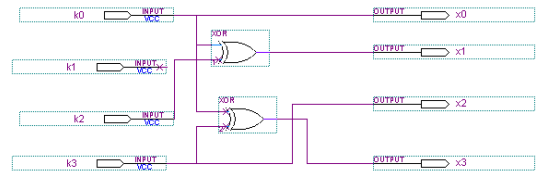


Fig. 2.  $K$  to  $Y_1$  converter logic circuit,  $Y_1=(x_0, x_1, x_2, x_3)$  and  $K=(k_0, k_1, k_2, k_3) \in GF(2^4)$   
그림 2.  $K$ 에서  $Y_1$ 으로 전환하는 논리회로

circuit for  $K$  to  $Y_1$  converter and Fig. 1 is  $K_1$  to  $K'_1$  Converter. Also  $(Y_1, Y_2)$  to real error position  $\in GF(2^8)$  can be found using  $GF(2^4)$  to  $GF(2^8)$  transform circuit<sup>[2]</sup>. In this section we saw how our new method becomes better than classical method. In next section we will compare the current method with the old method.

<Example 1>

If received polynomial is  $r(x)=\alpha^3x^2+\alpha^4x^5$  for  $c(x)=0$  (transmitted code polynomial) find 2 error locations and error values in  $GF(2^8)$ .

<sol>

$S_1=r(\alpha)=\alpha^{105}$   
 $S_2=r(\alpha^2)=\alpha^{119}$   
 $S_3=r(\alpha^3)=\alpha^{30}$   
 $S_4=r(\alpha^4)=\alpha^{110}$ , so from equation (2), we get  $\sigma_1 = \alpha^{225}$  and  $\sigma_2 = \alpha^7 \in GF(2^8)$ . Also  $K = \sigma_2 / (\sigma_1^2) = \alpha^{67}$ . Now transform into  $GF(2^4)$ . Then  $K = K_1 + \beta K_2$ ,  $(K_1, K_2) \in GF(2^4)$ . We get  $K_1 = \alpha^7$ ,  $K_2 = \alpha^5$ . From table 3,  $Y_2 = \alpha^7, \alpha^{13}$ . One of these  $Y_2$  is invalid. Let's assume  $Y_2 = \alpha^7$ . From (20)  $K'_1 = K_1 + \gamma Y_2^2 = \alpha^7 + \gamma \alpha^{14} = (1110) + \gamma (0011) = (1110) + (1000) = (0110) = \alpha^{13}$  from table 1. So from table 3, we see  $Y_1 = \alpha$  or  $\alpha^{12}$  (Notice  $\alpha^{12} + 1 = \alpha$ ). Therefore 2 Normalized solutions  $X_1$  is

$$X_1 = Y_1 + \beta Y_2 = (\alpha + \beta \alpha^7) \text{ or } (\alpha^{12} + \beta \alpha^7). \quad (18)$$

Transform (18) into  $GF(2^8)[2]$ , then

$$X_1 = \alpha^{35} \text{ or } \alpha^{32} \in GF(2^8). \quad (19)$$

So real solution to the equation (err locator polynomial) (1) is  $x = \sigma_1 X_1 = \alpha^{225} \cdot (\alpha^{35} \text{ or } \alpha^{32}) = \alpha^2$  or  $\alpha^5$  ( $\because \alpha^{255} = 1$ ). These are 2 correct error positions. If we assume  $Y_2 = \alpha^{13}$ ,  $K'_1 = \alpha^7 + \gamma (\alpha^{13})^2 = \alpha^7 + \gamma \alpha^{11}$

( $\because \alpha^{15} = 1$  for  $GF(2^4)$  elements) $=\alpha^7 + \alpha^{12} = \alpha^2 \in GF(2^4)$ .

From table 3 , there is no answer , so  $Y_2 = \alpha^{13}$  is invaild. Now 2 error values  $E_2$  and  $E_5$  are as follows. From equations (7) and (8)

$$E_2 = (\alpha^5 S_0 + S_1) / \alpha_1$$

$$= (\alpha^5 (\alpha^3 + \alpha^4) + \alpha^{105}) / \alpha^{225}$$

$$= \alpha^3 \in GF(2^8). \text{ and}$$

$$E_5 = S_0 + E_2 = \alpha^3 + \alpha^4 + \alpha^3$$

$$= \alpha^4 \in GF(2^8). \text{ These are correct !.}$$

<Example 2>

K to  $Y_1$  converter :

If  $K = \alpha^{13}$ , find  $Y_1(Y_1^2 + Y_1 = K_1)$ .

<Sol>

From equation (17),

$$K = [k_0, k_1, k_2, k_3] = [0, 1, 1, 0]$$

$$Y_1 = (y_{10}, y_{11}, y_{12}, y_{13})$$

$$= (k_0, k_0 + k_2, k_3, k_3 + k_0)$$

$$= (0, 1, 0, 0) = \alpha \text{ and another solution is}$$

$$Y_1 + 1 = (1, 1, 0, 0) = \alpha^{12}.$$

This is correct, when we see the ROM table 3.

#### IV. Comparison Between the New Method and Classical Method.

There are 2 big advantages over classical method, when we use new method to decode 2 error correcting RS code. First is ROM size reduction. When we compare table 2 and 3, we find 1/16 ROM size reduction. Second advantage is simpler logic for decoding circuit. For example ,to implement K to normalized error position converter, for classical method,we need 12 EXOR gates<sup>[7]</sup> , but

Table 4. Comparisons between Classical and New Methods.

표 4. 기존방식과 새방법의 비교표

	Table ROM size to find 2 error positions	Number of EXOR gates to Convert K to normalized error positions
Classical Method	(8bits/word) X (128 words)	12
New Method	(4bits/word) X (16 words)	2

Table 5. Comparisons between  $GF(2^4)$  Multiplier and  $GF(2^8)$  Multiplier.

표 5.  $GF(2^4)$ 승산기와  $GF(2^8)$ 승산기의 비교

	Number of EXOR Gates	Number of AND gates
$GF(2^4)$ case	15	16
$GF(2^8)$ case	73	64

for new method, we need just 2 EXOR gates as shown in equation 17. We summarizes the comparisons in table 4.

We see this reduction of logic gate counts not only makes the cost of the decoding circuit more economical but also the speed to find the error locations becomes greatly faster because of reduced propagation delay path.

#### V. Conclusions

Here we proposed new method to optimize 2error correcting Reed-Solomon(RS) decoder circuit. HDTV, Digital Audio device, Digital communication devices use the RS decoder to prote ct the digital data in the devices<sup>[3,7]</sup>. Norm ally the RS decoder in the devices corrects 2 ~ 4 symbol errors for 1 codeword. 4error correcting RS decoder is actually same as 2 error correcting RS decoder, if code word leng th becomes as half as its original word leng th. So 2 error correcting RS decoder circuit is very Important for the digital devices. In future we will try to optimize the 3 error correcting RS decoder and the circuit to determine the number of error symbols in one code word of RS decoder circuit.

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